

SOME METRICAL PROBLEMS ON S_n^*

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We present some problems related to right- (or bi-)invariant metrics d on the symmetric group of permutations S_n . Characterizations and constructions of bi-invariant extremal (in the corresponding convex cone) metrics are given, esp. for $n \leq 5$. We also consider special subspaces of the metric space (S_n, d) : unit balls, sets with prescribed distances (L -cliques), “hamiltonian” sets.

Here we give (for proofs, see [3]) some results and problems arising by analogy with extremal set systems. Related problems of coding type (with Hamming metric) are considered in [1, 4].

1. Invariance

S_n is the symmetric group of degree n whose elements (permutations) are denoted $\alpha, \beta, \gamma, \dots$, with 1 being the identity. It is endowed with an integer metric $d: S_n \times S_n \rightarrow \mathbb{N}$ which will always be in this paper *right-invariant*, i.e. $\forall \alpha, \beta, \gamma \ d(\alpha\gamma, \beta\gamma) = d(\alpha, \beta)$. Then $d(\alpha, \beta) = d(\alpha\beta^{-1}, 1) = p_d(\alpha\beta^{-1})$ is the *weight* of $\alpha\beta^{-1}$. We will denote $d(1, \alpha)$ by $d(\alpha)$.

Examples. $H(\sigma) = |\{i: \sigma(i) \neq i\}|$; $L_1(\sigma) = \sum |\sigma(i) - i|$; $L_\infty(\sigma) = \text{Max } |\sigma(i) - i|$; $T(\sigma) = \text{min. number of transpositions } t_i \text{ such that } t_i t_j \cdots \sigma = 1$. If d is also *left-invariant*, i.e. $\forall \alpha, \beta, \gamma \ d(\gamma\alpha, \gamma\beta) = d(\alpha, \beta)$, it is said to be *bi-invariant*.

Proposition 1. *The bi-invariance of d is equivalent to any of the following conditions:*

- (1) $\forall \alpha, \beta \ d(\alpha\beta) = d(\beta\alpha)$,
- (2) $\forall \alpha, \beta \ d(\alpha, \beta) = d(\alpha^{-1}, \beta^{-1})$,
- (3) $\forall \alpha, \beta \ d(\alpha) = d(\beta\alpha\beta^{-1})$.

Calling $N_i(\sigma)$ the number of cycles of length i in σ , with $\sum iN_i(\sigma) = n$, and $\mathcal{E}_\sigma$ the conjugacy class of σ , i.e. $\mathcal{E}_\sigma = \{\alpha\sigma\alpha^{-1}, \alpha \in S_n\}$ we get

Corollary. *H and T are bi-invariant.*

* Ce travail a été effectué pendant le séjour des auteurs au Centre de Recherche de Mathématiques, Appliquées, Université de Montréal.

It comes from the fact that $H(\sigma) = n - N_1(\sigma)$, $T(\sigma) = n - \sum N_i(\sigma)$ (Cayley) and (3) of Proposition 1.

2. Graphic distance

The distance d is *graphic* if $d(\alpha, \beta)$ is the length of the shortest path joining α and β in the graph whose vertex set is S_n and edge-set $\{(\sigma, \theta) : d(\sigma, \theta) = 1\}$. This is equivalent to saying: $\forall \alpha, \beta \ d(\alpha, \beta) \geq 2 \Rightarrow \exists \gamma$ between α and β , i.e. such that $d(\alpha, \beta) = d(\alpha, \gamma) + d(\gamma, \beta)$ (see [10]).

Proposition 2. *If d is a graphic distance (connected graph), then $E_d = \{\alpha : d(\alpha) = 1\}$ generates S_n . Reciprocally, if E is a symmetric set (i.e. $e \in E \Rightarrow e^{-1} \in E$), then the distance d_E (defined by $d_E(1) = 0$, $d_E(e) = 1$ ($\forall e \in E$), $d_E(\alpha)$ is the smallest number of $e_i \in E$, such that $\alpha = e_i e_j \dots$, $d_E(\alpha) = \infty$ when such a writing is impossible) is a graph weight, finite when E generates S_n .*

Proposition 3. *d_E is bi-invariant iff E is stable by conjugacy, i.e. $E = \bigcup \mathcal{E}_{\alpha}$.*

Some constructions. Let d_{E_1, E_2} be graph distances as in Proposition 2. Denote:

- (1) $d_{E_1} \wedge d_{E_2} = d_{E_1 \cup E_2}$,
- (2) $d_{E_1} \vee d_{E_2} = d_{E_1 \wedge E_2}$,
- (3) $d_{E_1} \circ d_{E_2} = d_{E_1 \circ E_2}$, where $E_1 \circ E_2 = \{\sigma_1 \sigma_2, \sigma_1 \in E_1, \sigma_2 \in E_2\}$.

They are graph weights, bi-invariant if d_{E_1} and d_{E_2} are bi-invariant. This is a partial answer to a question of [5].

3. Extremal bi-invariant metrics

Bi-invariant metrics form a convex cone over $\mathbb{R}$.

Proposition 4. *If E is exactly one conjugacy class, i.e. $E = \mathcal{E}_{\alpha}$, then d_E is an extreme ray of the cone of bi-invariant metrics.*

Examples. $E = \mathcal{E}_{(12)}$, the set of all transpositions; then $d_E = T$ is extremal. Let C_i denote a cycle of length i , $E_i = \{C_i\}$ generates S_n for even i . The associated distance $d_{\{C_i\}}$ is extremal. We have

$$H = T + \bigwedge_{i=2}^{\infty} d_{\{C_i\}}.$$

Thus H is not an extreme ray.

Proposition 5. Let $E_1 = \mathcal{E}_\alpha$, $E_2 = \mathcal{E}_\beta$, then d_{E_1} and d_{E_2} are Lipschitz-isomorph, i.e.

$$k_1^{-1}d_{E_1} \leq d_{E_2} \leq k_2d_{E_1}$$

with $k_1 = P_{E_1}(\beta)$, $k_2 = P_{E_2}(\alpha)$ independent of n .

The characterization of extremal bi-invariant metrics (graphic or not) seems a difficult problem.

Examples. A few bi-invariant metrics on S_5 (see Table 1).

Table 1

d	Cycle structure					
	(1) (2) (3) (45)	(1) (2) (345)	(1) (2345)	(1) (23) (45)	(12345) (12) (345)	
H	2	3	4	4	5	5
T	1	2	3	2	4	3
$d_{\{C_4\}}$	3	2	1	2	2	3
$T \wedge d_{\{C_3\}}$	1	1	2	2	2	2
$\bigwedge_{i=2}^{\infty} d_{\{C_i\}}$	1	1	1	2	1	2

4. Some special subspaces of the metric space (S_n, d)

4.1. Hamiltonian sets

Let E be a set of transformations, then $[12] (S_n, d_E)$ has a hamiltonian circuit if the graph G_E with vertex set $\{1, 2, \dots, n\}$ and edge set $\{(i, j): \exists t \in E, t(i) = j, t(j) = i\}$ is connected.

Examples. $G_E = K_n$, then $d_E = T$.

G_E is the path of length n (i.e. $1, 2, \dots, n$), then d_E is noted I .

G_E is the star with center 1, then d_E is noted U .

From $L_\infty \leq I$, one deduces that (S_n, L_∞) is also hamiltonian.

Proposition 6. If (S_n, d_{E_1}) and (S_n, d_{E_2}) are hamiltonian, then $(A_n, d_{E_1} \circ d_{E_2})$ is hamiltonian.

Example. $(A_n, U \circ U)$ is hamiltonian, with

$$U \circ U = d_E, \text{ where } E = \{t_1 \circ t_j: t_i = (1i), t_j = (1j)\}.$$

That is, one can generate A_n by performing permutations which are cycles of length 3 of the form $(ji1)$. Hence $(A_n, d_{\{C_3\}})$ is hamiltonian.

In [9] a similar question is investigated. Ring n bells in a “good” way, which may be restated as follows: find a hamiltonian circuit in (S_n, L_∞) with the extra condition that $\varphi(\sigma_i \sigma_{i+1}^{-1}) \cap \varphi(\sigma_{i+1} \sigma_{i+2}^{-1}) = \emptyset$, where $\varphi(\alpha)$ is the set of fixed points of α , and $\sigma_i, \sigma_{i+1}, \sigma_{i+2}$ are any three consecutive nodes of the circuit.

4.2. Metric basis and symmetries

A d -metric basis B is a set $B \subset S_n$ such that $\forall \sigma, \pi \in S_n$, if $\forall \beta \in B \quad d(\sigma, \beta) = d(\pi, \beta)$, then $\sigma = \pi$.

It is shown in [5] that $1 \cup \{C_2\} \cup \{C_3\}$ is a H -metric basis and this is used to prove that G_H , the group of isometries of (S_n, H) has order $2(n!)^2$ and contains as normal subgroup $F = \{f_{\sigma\pi}: f_{\sigma\pi}(\alpha) = \sigma\alpha\pi^{-1}\}$. For any bi-invariant d , it is easy to see that $G_d \supset G_H$. It would be interesting to find minimal d -metric basis for other d 's.

4.3. Spheres and balls

Let $S(d, n, r) = \{\sigma \in S_n: d(\sigma) = r\}$, $B(d, n, r) = \{\sigma \in S_n: d(\sigma) \leq r\}$ then

$$|S(H, n, r)| = \binom{n}{r} r! \sum_{i=0}^r \frac{(-1)^i}{i!} \simeq e^{-1} \binom{n}{r} r!,$$

$$|S(T, n, r)| = \sum_{\substack{(t_1 \dots t_n) \in \{1, 2, \dots, n\}^n \\ \sum t_i = n-r}} \frac{n!}{1^{t_1} t_1! \dots n^{t_n} t_n!}, \quad (\text{see [2]})$$

$$|B(L_\infty, n, 1)| = |B(L_\infty, n-1, 1)| + |B(L_\infty, n-2, 1)|,$$

$$|B(L_\infty, n, 2)| = 2 |B(L_\infty, n-1, 2)| + 2 |B(L_\infty, n-3, 2)| - |B(L_\infty, n-5, 2)|,$$

$$|S(I, n, r)| = \sum_{i=0}^{n-1} |S(I, n-1, r-i)|. \quad (\text{see [11]})$$

5. L -cliques in (S_n, d)

Let d be a right-invariant distance on S_n . Let L be a subset of $\{1, 2, \dots, n\}$. We call a subset A of S_n a L -clique (denote $\mathcal{A}(L)$) if $d(x, y) \in L$ for any $x, y \in A$, $x \neq y$. We call $\mathcal{A}(L)$ l -code if $L = \{l, l+1, \dots, n\}$, l -anticode if $L = \{1, 2, \dots, l-1\}$, l -equidistant code if $L = \{l\}$ for some integer l . Let us fix a L -clique A .

For any subset $S \subseteq S_n$ denote $\mathcal{A}_S(L)$ any subset $B \subseteq S$ such that B is a L -clique.

Proposition 7. Let $S \subseteq S_n$, then

$$|A| \leq \max |\mathcal{A}_S(L)| \cdot \frac{n!}{|S|}$$

if either d is bi-invariant or A is symmetric.

Proposition 7 follows from *density bound*

$$\frac{|A|}{|S_n|} \leq \frac{|\mathcal{A}_S(L)|}{|S|}.$$

Let $q(\alpha, \beta)$ be a right-invariant function $q: S_n \times S_n \rightarrow \mathbb{R}$ such that

- (a) matrix $(q(\alpha, \beta))$ of order $n!$ has only nonnegative eigenvalues and
- (b) $q(\alpha, \beta) \leq 0$ whenever $d(\alpha, \beta) \in L$.

Proposition 8.

$$|A| \leq (n!)^2 \frac{\max_{\alpha \in S_n} q(\alpha, \alpha)}{\sum_{\alpha, \beta \in S_n} q(\alpha, \beta)}.$$

Proposition 8 follows from *averaging bound* [8]

$$\frac{\sum_{\alpha, \beta \in S_n} q(\alpha, \beta)}{|S_n|^2} \leq \frac{\sum_{\alpha, \beta \in A} q(\alpha, \beta)}{|A|^2}$$

Denote $\bar{L} = \{1, 2, \dots, n\} - L$.

Proposition 9.

$$|A| \leq \frac{n!}{\max |\mathcal{A}(\bar{L})|}$$

if either d is bi-invariant or A is symmetric.

Proposition 9 follows from *duality bound* [7]

$$|\mathcal{A}(L)| \cdot |\mathcal{A}(\bar{L})| \leq |S_n|.$$

We give now two applications of Proposition 9 for the case $L = \{l, l+1, \dots, n\}$, i.e. L -clique A is a l -code.

Denote $A_1 = B(d, n, \lfloor \frac{1}{2}(l-1) \rfloor)$. Of course, A_1 is a symmetric l -anticode. Denote A_2 the stabilizer of a smallest subset $M \subseteq \{1, 2, \dots, n\}$ such that this stabilizer is a $\bar{L}$ -clique (i.e. l -anticode); of course, A_2 is symmetric.

Corollary. For $L = \{l+1, l+2, \dots, n\}$

- (i) $|A| \leq n!/|A_1|$,
- (ii) $|A| \leq n!/|A_2|$.

Explicit values for some $|A_1|$ are derivable from Section 4.3. For (ii) one obviously has $|A_2| = (n - |M|)!$ but a general relation between $|M|$ and l is complicated (see [2]).

Two well-known upper bounds for codes in $(\mathbb{F}_q)^n$ with Hamming distance (Hamming–Rao and Singleton’s bounds) corresponds to (i), (ii). For $d = H$: we have equality in (ii) iff A is a sharply $(n - l + 1)$ -transitive subset of S_n ; equality in (i) will give the analog of perfect codes.

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